

# Role Swap: When the Follower Leads and the Leader Follows

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## ONLINE APPENDIX

### APPENDIX A. PROOF OF PROPOSITION 1 FOR THE 3×3 GAME

The steps to derive the three types of equilibrium conditions for the 3×3 game, as well as the intuition, are analogous to those for the 2×2 game appearing in the main text. This is apparent when comparing Figure 4 below with Figure 2 in the main text. It is therefore unnecessary to provide detailed explanations here.

We will start again with the conditions for  $M$ 's Dominance region, i.e. for  $(M_1^S M_2^S, F_1^S F_2^S)$  to be the unique subgame perfect equilibrium based on elimination of strictly dominated strategies.  $F$ 's Yielding condition is composed of four inequalities, which can be compactly written as follows.

$$(28) \quad \begin{aligned} & \underbrace{(1-m)t + \overbrace{my}^{F_2^S}}_{M_2^S} > \max \left\{ \underbrace{(1-m)v + \overbrace{mr}^{F_2^B}}_{M_2^S}, \underbrace{(1-m)x + \overbrace{mw}^{F_2^C}}_{M_2^S} \right\}, \\ & \underbrace{(1-m)t + \overbrace{mz}^{F_2^S}}_{M_2^S} > \max \left\{ \underbrace{(1-m)v + \overbrace{mu}^{F_2^B}}_{M_2^S}, \underbrace{(1-m)x + \overbrace{ms}^{F_2^C}}_{M_2^S} \right\}. \end{aligned}$$

Switch to  $M_2^B$       Switch to  $M_2^C$       Switch to  $M_2^B$       Switch to  $M_2^C$

Combining these conditions and rearranging implies

$$(29) \quad m < \bar{m}_{3 \times 3} = \min \left\{ \frac{t-v}{t-v+r-y}, \frac{t-v}{t-v+u-z}, \frac{t-x}{t-x+s-z} \right\} \stackrel{(3)}{=} \frac{2}{7}.$$

Assuming (29) holds and moving backwards,  $M$ 's Sticking condition is similarly composed of four constraints

$$(30) \quad \begin{aligned} & \underbrace{(1-f)e + \overbrace{fa}^{M_2^S}}_{F_2^B} > \max \left\{ \underbrace{(1-f)c + \overbrace{fi}^{M_2^B}}_{F_2^B}, \underbrace{(1-f)h + \overbrace{fg}^{M_2^C}}_{F_2^B} \right\}, \\ & \underbrace{(1-f)d + \overbrace{fa}^{M_2^S}}_{F_2^C} > \max \left\{ \underbrace{(1-f)j + \overbrace{fi}^{M_2^B}}_{F_2^C}, \underbrace{(1-f)b + \overbrace{fg}^{M_2^C}}_{F_2^C} \right\}. \end{aligned}$$

switch to  $F_2^S$       switch to  $F_2^S$       switch to  $F_2^S$       switch to  $F_2^S$

They can be rearranged into

$$(31) \quad f > \max \left\{ \frac{c-e}{a-i+c-e}, \frac{b-d}{a-g+b-d} \right\} \stackrel{(3)}{=} \frac{3}{8}.$$

Finally, assuming (31) is satisfied,  $M$ 's Contest condition can be expressed as

$$(32) \quad \begin{array}{c} \overbrace{(1-f)e + fa}^{M_1^S M_2^S} > \underbrace{c}_{M_1^B M_2^B}, \\ \underbrace{F_1^B \text{ as can't revise}} \quad \text{switch to } F_2^S \quad \underbrace{F_1^B} \\ \overbrace{(1-f)d + fa}^{M_1^S M_2^S} > \underbrace{b}_{M_1^C M_2^C}. \\ \underbrace{F_1^C \text{ as can't revise}} \quad \text{switch to } F_2^S \quad \underbrace{F_1^C} \end{array}$$

After some manipulations we obtain

$$(33) \quad f > \hat{f}_{3 \times 3} = \max \left\{ \frac{c-e}{a-e}, \frac{b-d}{a-d} \right\} \stackrel{(3)}{=} \frac{3}{5}.$$

It is apparent that like in the  $2 \times 2$  game reported in the main text, in the  $3 \times 3$  game the Contest condition in (33) is stronger than the Sticking condition in (31) for all general parameter values. This means that the necessary and sufficient conditions for  $M$ 's Dominance region are jointly (29) and (33).

The analogous circumstances for  $F$ 's Dominance region, namely  $M$ 's Yielding and  $F$ 's Contest conditions, can again be derived by symmetry

$$(34) \quad \begin{array}{l} f < \bar{f}_{3 \times 3} = \min \left\{ \frac{c-e}{c-e+a-i}, \frac{c-e}{c-e+d-j}, \frac{c-h}{c-h+b-j} \right\} \stackrel{(3)}{=} \frac{2}{7}, \text{ and} \\ m > \hat{m}_{3 \times 3} = \max \left\{ \frac{t-v}{r-v}, \frac{s-u}{r-u} \right\} \stackrel{(3)}{=} \frac{3}{5}. \end{array}$$

This completes the proof of Proposition 1 for the  $3 \times 3$  game.

## APPENDIX B. PROOF OF PROPOSITION 2 FOR THE $3 \times 3$ GAME

Using the same logic as in the  $2 \times 2$  Battle of the sexes, in the  $3 \times 3$  version the condition for  $M$  to secure his preferred outcome despite being the Stochastic follower is

$$\hat{f}_{3 \times 3} = \max \left\{ \frac{c-e}{a-e}, \frac{b-d}{a-d} \right\} < \bar{m}_{3 \times 3} = \min \left\{ \frac{t-v}{t-v+r-y}, \frac{t-v}{t-v+u-z}, \frac{t-x}{t-x+s-z} \right\}.$$

Analogously, for  $F$  this occurs if

$$\hat{m}_{3 \times 3} = \max \left\{ \frac{t-v}{r-v}, \frac{s-u}{r-u} \right\} < \bar{f}_{3 \times 3} = \min \left\{ \frac{c-e}{c-e+a-i}, \frac{c-e}{c-e+d-j}, \frac{c-h}{c-h+b-j} \right\}.$$

To offer a numerical example, consider the specific payoffs in (3) and change the female's  $(B, B)$  payoff from  $r = 5$  to  $r = 11$ . This implies  $\hat{m}_{3 \times 3} = \frac{3}{11}$ , which is smaller than  $\bar{f}_{3 \times 3} = \frac{2}{7}$ . This means that  $F$ 's Dominance region crosses the 45-degree line, so  $F$  can ensure its Stackelberg payoffs even from the position of the Stochastic follower.

### The 3x3 game

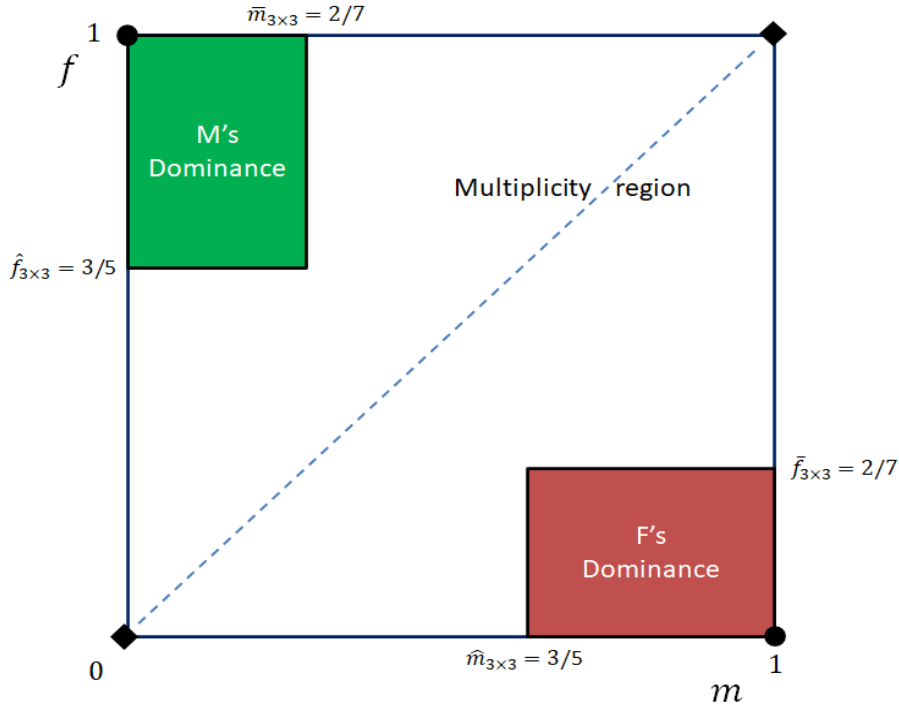


FIGURE 1. Equilibrium regions in the  $3 \times 3$  Battle of the sexes under Stochastic leadership, featuring revision probabilities  $m$  and  $f$ . The reported numerical thresholds apply for the specific payoffs in (3).

#### APPENDIX C. PROOF OF PROPOSITION 3 FOR THE $3 \times 3$ GAME

Let us now examine the  $3 \times 3$  Stag hunt, focusing again on  $M$ 's Dominance region.  $F$ 's Yielding condition is composed of the following

$$(35) \quad \underbrace{\overbrace{(1-m)\mathcal{R}}^{F_2^S} + \underbrace{m\mathcal{Z}}_{\text{switch to } M_2^H}}_{M_2^S \text{ as can't revise}} > \max \left\{ \underbrace{\overbrace{(1-m)\mathcal{T}}^{F_2^H} + \underbrace{m\mathcal{W}}_{\text{switch to } M_2^H}}_{M_2^S}, \underbrace{\overbrace{(1-m)\mathcal{S}}^{F_2^B} + \underbrace{m\mathcal{Y}}_{\text{switch to } M_2^H}}_{M_2^S} \right\},$$

$$\underbrace{\overbrace{(1-m)\mathcal{R}}^{F_2^S} + \underbrace{m\mathcal{X}}_{M_2^B}}_{M_2^S} > \max \left\{ \underbrace{\overbrace{(1-m)\mathcal{T}}^{F_2^H} + \underbrace{m\mathcal{V}}_{M_2^B}}_{M_2^S}, \underbrace{\overbrace{(1-m)\mathcal{S}}^{F_2^B} + \underbrace{m\mathcal{U}}_{M_2^B}}_{M_2^S} \right\}.$$

Rearranging and combining these inequalities implies

$$(36) \quad m < \min \left\{ \frac{\mathcal{R} - \mathcal{T}}{\mathcal{R} - \mathcal{T} + \mathcal{W} - \mathcal{Z}}, \frac{\mathcal{R} - \mathcal{S}}{\mathcal{R} - \mathcal{S} + \mathcal{Y} - \mathcal{Z}}, \frac{\mathcal{R} - \mathcal{T}}{\mathcal{R} - \mathcal{T} + \mathcal{V} - \mathcal{X}}, \frac{\mathcal{R} - \mathcal{S}}{\mathcal{R} - \mathcal{S} + \mathcal{U} - \mathcal{X}} \right\}.$$

Using (17) suggests that the first, second and fourth fractions in (36) can be the smallest, which means we can streamline it as

$$(37) \quad m < \bar{m}_{3 \times 3}^{SH} = \min \left\{ \frac{\mathcal{R} - \mathcal{T}}{\mathcal{R} - \mathcal{T} + \mathcal{W} - \mathcal{Z}}, \frac{\mathcal{R} - \mathcal{S}}{\mathcal{R} - \mathcal{S} + \mathcal{Y} - \mathcal{Z}}, \frac{\mathcal{R} - \mathcal{S}}{\mathcal{R} - \mathcal{S} + \mathcal{U} - \mathcal{X}} \right\} \stackrel{(15)}{=} \frac{1}{2}.$$

Moving backwards and assuming (37) is satisfied  $M$ 's Sticking condition can be written as

$$(38) \quad \underbrace{\overbrace{(1-f)\mathcal{J} + f\mathcal{A}}^{M_2^S}}_{F_2^H} > \max \left\{ \underbrace{\overbrace{(1-f)\mathcal{G} + f\mathcal{C}}^{M_2^H}}_{F_2^H}, \underbrace{\overbrace{(1-f)\mathcal{I} + f\mathcal{B}}^{M_2^B}}_{F_2^S} \right\},$$

$$\underbrace{\overbrace{(1-f)\mathcal{H} + f\mathcal{A}}^{M_2^S}}_{F_2^B} > \max \left\{ \underbrace{\overbrace{(1-f)\mathcal{E} + f\mathcal{C}}^{M_2^H}}_{F_2^B}, \underbrace{\overbrace{(1-f)\mathcal{D} + f\mathcal{B}}^{M_2^B}}_{F_2^S} \right\}.$$

Combining them implies

$$(39) \quad f > \max \left\{ \frac{\mathcal{G} - \mathcal{J}}{\mathcal{G} - \mathcal{J} + \mathcal{A} - \mathcal{C}}, \frac{\mathcal{I} - \mathcal{J}}{\mathcal{I} - \mathcal{J} + \mathcal{A} - \mathcal{B}}, \frac{\mathcal{E} - \mathcal{H}}{\mathcal{E} - \mathcal{H} + \mathcal{A} - \mathcal{C}}, \frac{\mathcal{D} - \mathcal{H}}{\mathcal{D} - \mathcal{H} + \mathcal{A} - \mathcal{B}} \right\}.$$

It is apparent from (17) that the first, second and fourth fractions in (39) can be binding, which allows us to reduce it to

$$(40) \quad f > \hat{f}_{3 \times 3}^{SH} = \max \left\{ \frac{\mathcal{G} - \mathcal{J}}{\mathcal{G} - \mathcal{J} + \mathcal{A} - \mathcal{C}}, \frac{\mathcal{I} - \mathcal{J}}{\mathcal{I} - \mathcal{J} + \mathcal{A} - \mathcal{B}}, \frac{\mathcal{D} - \mathcal{H}}{\mathcal{D} - \mathcal{H} + \mathcal{A} - \mathcal{B}} \right\} \stackrel{(15)}{=} \frac{1}{2}.$$

Assuming that both (37) and (40) hold  $M$ 's Contest condition is

$$(41) \quad \underbrace{\overbrace{(1-f)\mathcal{J} + f\mathcal{A}}^{M_1^S M_2^S}}_{F_1^H} > \underbrace{\overbrace{\mathcal{G}}^{M_1^H M_2^H}} & \text{and} & \underbrace{\overbrace{(1-f)\mathcal{H} + f\mathcal{A}}^{M_1^S M_2^S}}_{F_1^B} > \underbrace{\overbrace{\mathcal{D}}^{M_1^B M_2^B}}.$$

Upon rearranging, we obtain

$$(42) \quad f > \max \left\{ \frac{\mathcal{G} - \mathcal{J}}{\mathcal{A} - \mathcal{J}}, \frac{\mathcal{D} - \mathcal{H}}{\mathcal{A} - \mathcal{H}} \right\} \stackrel{(15)}{=} \frac{1}{3}.$$

The Contest condition is always weaker than the Sticking condition, as was the case in the  $2 \times 2$  Stag hunt. As such, the necessary and sufficient conditions in the  $3 \times 3$  game for  $M$ 's dominance region are (37) and (40).

Similarly, the necessary and sufficient conditions for  $F$ 's dominance region are

$$f < \bar{f}_{3 \times 3}^{SH} = \min \left\{ \frac{\mathcal{A} - \mathcal{C}}{\mathcal{A} - \mathcal{C} + \mathcal{G} - \mathcal{J}}, \frac{\mathcal{A} - \mathcal{B}}{\mathcal{A} - \mathcal{B} + \mathcal{I} - \mathcal{J}}, \frac{\mathcal{A} - \mathcal{B}}{\mathcal{A} - \mathcal{B} + \mathcal{D} - \mathcal{H}} \right\} \stackrel{(15)}{=} \frac{1}{2},$$

and

$$m > \hat{m}_{3 \times 3}^{SH} = \max \left\{ \frac{\mathcal{W} - \mathcal{Z}}{\mathcal{W} - \mathcal{Z} + \mathcal{R} - \mathcal{T}}, \frac{\mathcal{Y} - \mathcal{Z}}{\mathcal{Y} - \mathcal{Z} + \mathcal{R} - \mathcal{S}}, \frac{\mathcal{U} - \mathcal{X}}{\mathcal{U} - \mathcal{X} + \mathcal{R} - \mathcal{S}} \right\} \stackrel{(15)}{=} \frac{1}{2}.$$

The condition for  $M$ 's Dominance region to cross the 45-degree line is

$$\begin{aligned} \hat{f}_{3 \times 3}^{SH} &= \max \left\{ \frac{G-J}{G-J+A-C}, \frac{I-J}{I-J+A-B}, \frac{D-H}{D-H+A-B} \right\} < \\ \bar{m}_{3 \times 3}^{SH} &= \min \left\{ \frac{R-T}{R-T+W-Z}, \frac{R-S}{R-S+Y-Z}, \frac{R-S}{R-S+U-X} \right\}, \end{aligned}$$

whereas the analogous one for  $F$ 's Dominance region is

$$\begin{aligned} \hat{m}_{3 \times 3}^{SH} &= \max \left\{ \frac{W-Z}{W-Z+R-T}, \frac{Y-Z}{Y-Z+R-S}, \frac{U-X}{U-X+R-S} \right\} < \\ \bar{f}_{3 \times 3}^{SH} &= \min \left\{ \frac{A-C}{A-C+G-J}, \frac{A-B}{A-B+I-J}, \frac{A-B}{A-B+D-H} \right\}. \end{aligned}$$

This completes the proof of Proposition 3 for the  $3 \times 3$  game.

#### APPENDIX D. PROOF OF PROPOSITION 4 FOR THE $3 \times 3$ GAME

Turning to the Hawk and dove game, let us again start with the conditions for  $M$ 's Dominance region in the  $3 \times 3$  game. The steps and intuition are again analogous to the  $2 \times 2$  game (as well as to the Battle of the sexes) above.  $F$ 's Yielding condition requires

$$(43) \quad \begin{aligned} \underbrace{(1-m)W + mU}_{M_2^H \quad M_2^D} &> \max \left\{ \underbrace{(1-m)Z + mR}_{M_2^H \quad M_2^D}, \underbrace{(1-m)Y + mS}_{M_2^H \quad M_2^D} \right\}, \\ \underbrace{(1-m)W + mV}_{M_2^H \quad M_2^O} &> \max \left\{ \underbrace{(1-m)Z + mX}_{M_2^H \quad M_2^O}, \underbrace{(1-m)Y + mT}_{M_2^H \quad M_2^O} \right\}, \end{aligned}$$

which can be rearranged into

$$(44) \quad m < \min \left\{ \frac{W-Z}{W-Z+R-U}, \frac{W-Y}{W-Y+S-U}, \frac{W-Z}{W-Z+X-V}, \frac{W-Y}{W-Y+T-V} \right\}.$$

Once we utilize (22) we can note that the first and fourth fractions in (44) can be the smallest of the four. The required condition can thus be simplified into

$$(45) \quad m < \bar{m}_{3 \times 3}^{HD} = \min \left\{ \frac{W-Z}{W-Z+R-U}, \frac{W-Y}{W-Y+T-V} \right\} \stackrel{(23)}{=} \frac{5}{8}.$$

Assuming (5) is satisfied,  $M$ 's Sticking condition requires

$$(46) \quad \begin{aligned} \underbrace{(1-f)J + fA}_{F_2^H \quad F_2^D} &> \max \left\{ \underbrace{(1-f)G + fD}_{F_2^H \quad F_2^D}, \underbrace{(1-f)I + fB}_{F_2^H \quad F_2^D} \right\}, \\ \underbrace{(1-f)H + fA}_{F_2^O \quad F_2^D} &> \max \left\{ \underbrace{(1-f)E + fD}_{F_2^O \quad F_2^D}, \underbrace{(1-f)C + fB}_{F_2^O \quad F_2^D} \right\}, \end{aligned}$$

which can be expressed as

$$f > \max \left\{ \frac{G-J}{G-J+A-D}, \frac{I-J}{I-J+A-B}, \frac{E-H}{E-H+A-D}, \frac{C-H}{C-H+A-B} \right\}.$$

Using (22), we can check that the first, second and forth fractions can be highest. So we can rewrite this condition as

$$(47) \quad f > \hat{f}_{3 \times 3}^{HD} = \max \left\{ \frac{G-J}{G-J+A-D}, \frac{I-J}{I-J+A-B}, \frac{C-H}{C-H+A-B} \right\} \stackrel{(23)}{=} \frac{5}{6}.$$

Moving backwards and assuming the Yielding and the Sticking conditions both hold,  $M$ 's Contest condition is as follows

$$(48) \quad \underbrace{(1-f)J + fA}_{\substack{M_1^H M_2^H \\ F_1^H}} > \underbrace{G}_{M_1^D M_2^D} \quad \text{and} \quad \underbrace{(1-f)H + fA}_{\substack{M_1^H M_2^H \\ F_1^O}} > \underbrace{C}_{M_1^O M_2^O}.$$

Rearranging this yields

$$(49) \quad f > \max \left\{ \frac{G-J}{A-J}, \frac{C-H}{A-H} \right\}.$$

It is clear that in the Hawk and dove game, like in the Battle of the sexes, the Contest condition is stronger than the Sticking condition in both the  $2 \times 2$  and  $3 \times 3$  games. Similarly, the conditions for  $F$ 's dominance region are:

$$f < \bar{f}_{3 \times 3}^{HD} = \min \left\{ \frac{G-J}{G-J+A-D}, \frac{G-I}{G-I+C-E} \right\} \stackrel{(23)}{=} \frac{5}{8},$$

and

$$m > \hat{m}_{3 \times 3}^{HD} = \max \left\{ \frac{W-Z}{W-Z+R-U}, \frac{Y-Z}{Y-Z+R-S}, \frac{T-X}{T-X+R-S} \right\} \stackrel{(23)}{=} \frac{5}{6}.$$

The condition for the role swap, in which  $M$  behaves as the Stackelberg leader despite being the Stochastic follower is

$$\hat{f}_{3 \times 3}^{HD} = \max \left\{ \frac{G-J}{G-J+A-D}, \frac{I-J}{I-J+A-B}, \frac{C-H}{C-H+A-B} \right\} < \bar{m}_{3 \times 3}^{HD} = \min \left\{ \frac{W-Z}{W-Z+R-U}, \frac{W-Y}{W-Y+T-V} \right\}$$

and for  $F$  the analogous conditions is

$$\hat{m}_{3 \times 3}^{HD} = \max \left\{ \frac{W-Z}{W-Z+R-U}, \frac{Y-Z}{Y-Z+R-S}, \frac{T-X}{T-X+R-S} \right\} < \bar{f}_{3 \times 3}^{HD} = \min \left\{ \frac{G-J}{G-J+A-D}, \frac{G-I}{G-I+C-E} \right\}$$

This completes the proof of Proposition 4 for the  $3 \times 3$  game.